

Global Optimization of Mixed-Integer Nonlinear Programs

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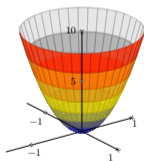


Introduction

Mixed-Integer Nonlinear Programs

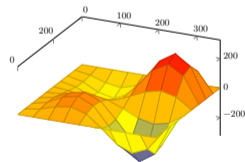
$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax \leq b \\ & g_k(x) \leq 0 \quad \forall k \in [m] \\ & x_i \in [\ell_i, u_i] \quad \forall i \in [n] \\ & x_i \in \mathbb{Z} \quad \forall i \in \mathcal{I} \subseteq [n] \end{aligned}$$

The nonlinear part: functions $g_k \in C^1([\ell, u], \mathbb{R})$:



convex

or



nonconvex

Examples of Nonlinearities

- Variable **multiplier** $p \in [0, 1]$ of variable quantity q : qp . Example: **water treatment** unit



- AC power flow** - nonlinear function of voltage magnitudes and angles



$$p_{ij} = g_{ij}V_i^2 - g_{ij}V_iV_j \cos(\theta_{ij}) + b_{ij}V_iV_j \sin(\theta_{ij})$$

- Distance** constraints



$$(x - x_0)^2 + (y - y_0)^2 \leq R$$

- etc.

Solving a Mixed-Integer Optimization Problem

Two major tasks:

1. Finding and improving feasible solutions (**primal side**)
 - **Ensure feasibility**, not necessarily maintaining optimality
 - Important for many practical applications
2. Proving optimality (**dual side**)
 - **Ensure optimality**, not necessarily maintaining feasibility
 - Necessary in order to actually solve the problem
 - Provides guarantees on solution quality

Linked by:

3. Strategy
 - **Ensure convergence**
 - Divide: branching, decompositions, ...
 - Put together all components to find a solution that is **feasible** and **optimal** (or within a proven gap from the optimum)

Nonlinearity Brings New Challenges

1. Primal side

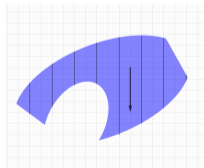
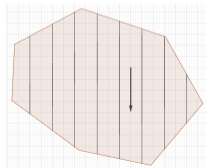
- Feasible solutions must also satisfy nonlinear constraints
- If nonconvex: local optima, local infeasibility

2. Dual side

- NLP relaxations capture the problem better, LP relaxations are faster
- Strong cuts needed for various nonlinearities
- If nonconvex: straightforward continuous relaxation no longer provides a lower bound

3. Strategy

- Need to account for all of the above
- Warmstart for NLP is less efficient than for LP
- More numerical issues
- NLP solving is less efficient and reliable than LP



Finding Feasible Solutions

Primal Heuristics

The goal of primal heuristics is to **find solutions** that are:

- **feasible** (satisfying all constraints) and
- **good quality** (solutions with lower objective value are preferable).

The best solution found so far is referred to as **best feasible** or **incumbent**. It provides an **upper/primal bound** on the optimal value.

Common theme in primal heuristics: **restrict the problem** to obtain an 'easier' subproblem for which a feasible solution can be found.

Nonconvex case: NLP subproblems are usually solved to local optimality.

- Local optima are still feasible solutions
- Not finding the global optimum affects the quality of upper bounds

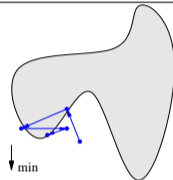
Primal Heuristics for MINLPs

MILP heuristics

- Can be applied to MINLPs (solutions violating nonlinear constraints can be passed to NLP local search).
-

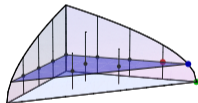
NLP local search

- Fix integer variables** to values at reference point; solve the NLP.
- Reference point examples: integer feasible solution of the LP relaxation, solution from an MILP heuristic.



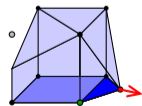
Undercover

- Fix some variables** so that constraints become **linear**; solve the MILP.



Sub-MINLP

- Extensions of MILP large neighbourhood search heuristics.
- Search around **promising** solutions.
- The region is restricted by additional constraints and/or fixing variables.



Proving Optimality

Proving Optimality

- Using relaxations for finding lower bounds
- Relaxations for convex MINLPs
- Relaxations for nonconvex MINLPs
- Managing cuts: initial cuts and dynamically added cuts
- How to strengthen relaxations

Finding Lower Bounds: Relaxations

A **relaxation** R of a **feasible set** F is a set such that $F \subseteq R$.

Requirement: the relaxed problem should be **efficiently** solvable to **global** optimality.

Relaxations can be:

- **Convex:** NLP solutions are globally optimal, infeasibility detection is reliable
- **Linear:** solving is more efficient, good for warmstarting

MILP and MINLP relaxations are sometimes used as well.

It is preferable that relaxations:

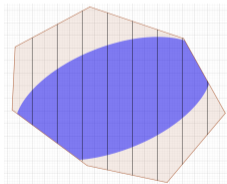
- Are **tight:** small $R \setminus F$, dual bound close to the optimal value
- Are **compact:** avoid excessive numbers of constraints and variables
- Have **reasonable numerics**

Relaxations for Convex MINLPs

- Relax integrality



- Replace the nonlinear set with a linear outer approximation



- Relax integrality + linear outer approximation \rightarrow LP relaxation

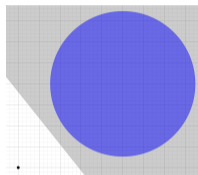
Outer Approximating Convex Constraints

A **linear inequality** $ax \leq b$ is **valid** if $x \in F \Rightarrow ax \leq b$ (**cutting planes**, or **cuts**, are valid inequalities)

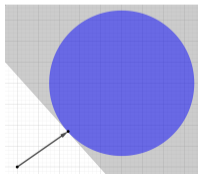
Given constraint $g(x) \leq 0$ and a **reference point** \hat{x} , one can build:

Gradient cuts (Kelley):

$$g(\hat{x}) + \nabla g(\hat{x})(x - \hat{x}) \leq 0$$



Projected cuts: same, but move \hat{x} to the boundary of F



Convex Relaxations for Nonconvex MINLPs

Relaxing integrality no longer provides a lower bound, and gradient cuts are no longer valid \Rightarrow construct a convex relaxation.

The best relaxation is $\text{conv}(F)$: **convex hull** of F , i.e. the smallest convex set containing F . In general, **cannot be constructed explicitly**.

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- Find **convex underestimators** g_k^{cv} of functions g_k : $g_k^{cv}(x) \leq g_k(x) \quad \forall x \in [l, u]$

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- Find and combine relaxations of **simple functions**

Examples of **simple functions**: x^2 , x^k , \sqrt{x} , xy , etc.

Combining Relaxations

Expression trees

Algebraic structure of nonlinear constraints can be stored in **one** directed acyclic graph:

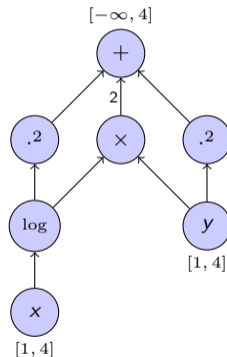
- nodes: variables, operations, constraints
- arcs: flow of computation

Underestimators of compositions

Find underestimator for $g(x) = \phi(\psi_1(x), \dots, \psi_p(x))$, where functions can be convexified directly.

- McCormick relaxations** for factorable functions: piecewise continuous relaxations utilising convex and concave envelopes of ϕ and ψ_j .
- Auxiliary variable** method: introduce variables $y_j = \psi_j(x)$. Then $g(x) = \phi(y_1, \dots, y_p)$. Enables individual handling of each function.

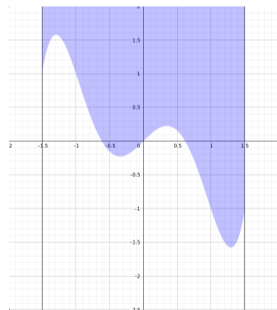
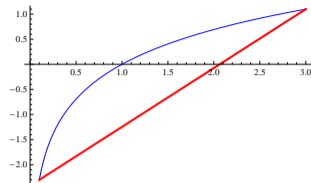
$$\log(x)^2 + 2 \log(x)y + y^2$$



Linear Relaxations for Nonconvex MINLPs

- If possible, **directly** construct linear underestimators for nonconvex functions
 - Secants for concave functions
 - McCormick envelopes for bilinear products
 - etc.

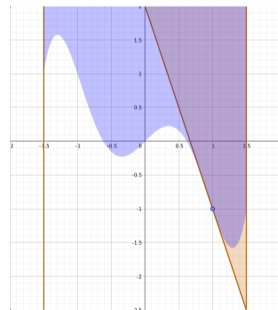
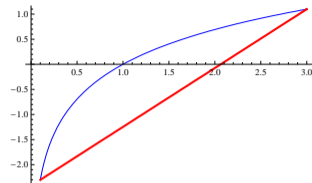
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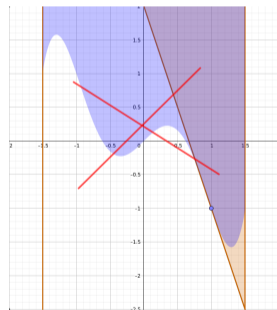
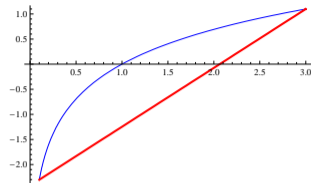
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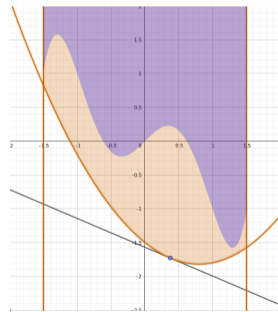
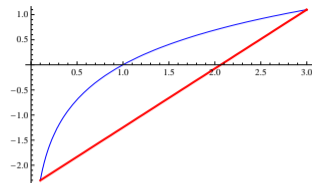
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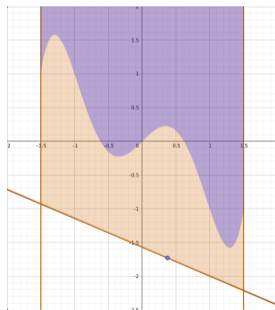
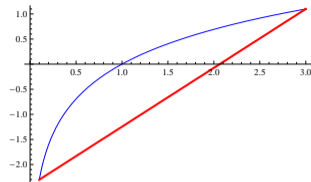
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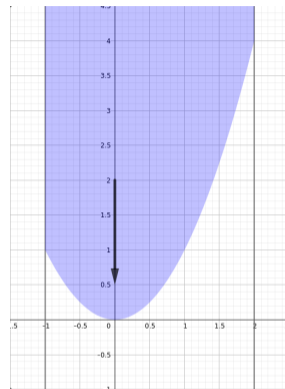
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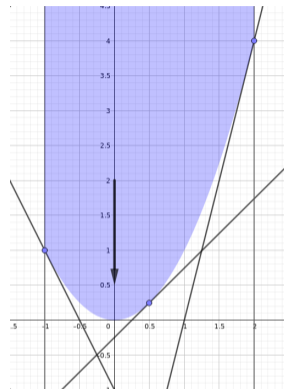
Cut Generation



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Initial cuts

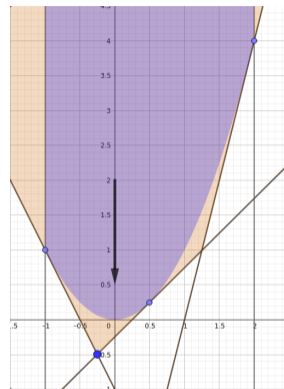
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- Reference points chosen based on feasible set only
- Aiming for a compact formulation that roughly captures F and yields a reference point for separation



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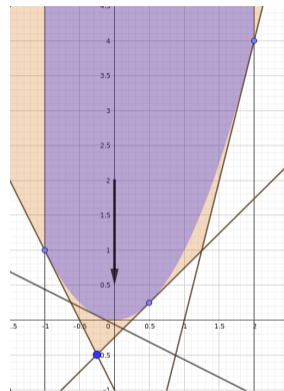
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- Valid inequalities $ax \leq b$ violated by \hat{x} : $a\hat{x} > b$
- Thus \hat{x} is **separated** from F



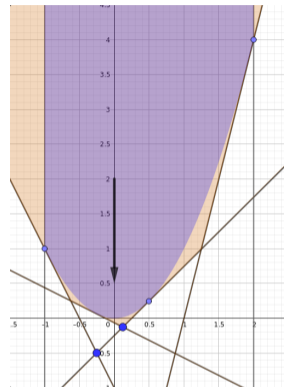
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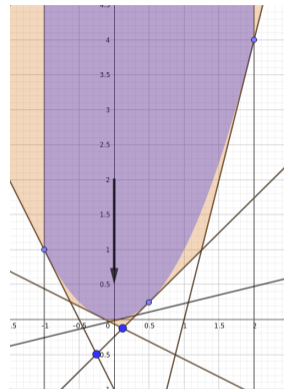
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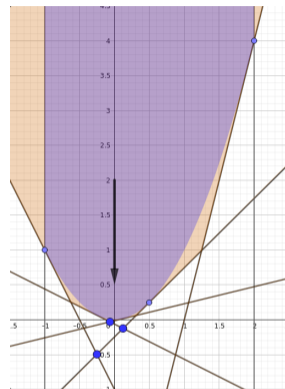
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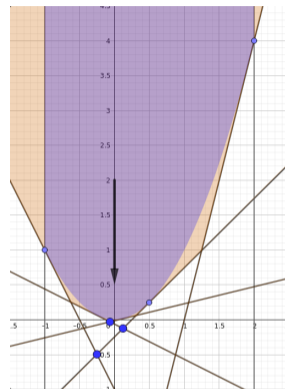
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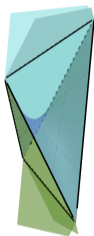
Cut selection: choose from violated cuts using various criteria for cut usefulness (e.g. violation, orthogonality, density, etc.).

Strengthening Relaxations: Tighter Variable Bounds

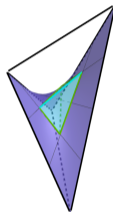
Tighter bounds \Rightarrow tighter relaxations.

Example: McCormick relaxation of a bilinear product relation $z = xy$:

$$\begin{aligned}z &\leq x^u y + xy^l - x^u y^l \\z &\leq x^l y + xy^u - x^l y^u \\z &\geq x^l y + xy^l - x^l y^l \\z &\geq x^u y + xy^u - x^u y^u\end{aligned}$$



$$(x, y) \in [-1, 2] \times [-2, 2]$$



$$(x, y) \in [0, 1] \times [-1, 1]$$

Tighter bounds obtained from:

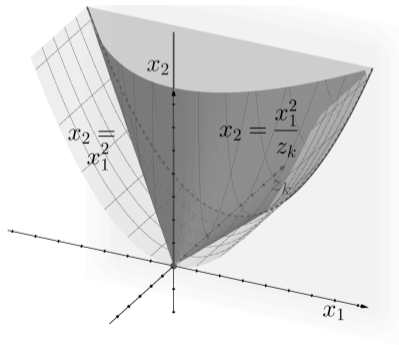
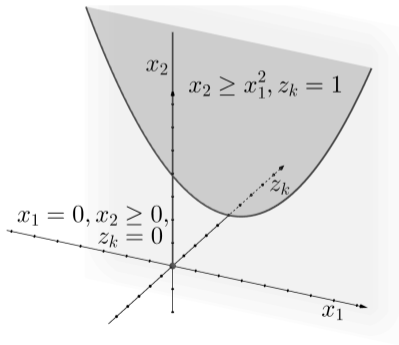
- **Branching** (more on this in the Strategy section)
- Specialized **bound tightening** techniques
- **Piecewise-linear** relaxations

Strengthening Relaxations: Using More Constraints

More constraints \Rightarrow tighter relaxations.

Example: **perspective formulations**. Use an additional constraint that requires x to be **semicontinuous**.

$$g(x) \leq 0, \quad lz \leq x \leq uz$$

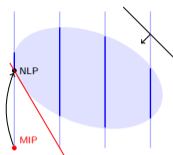


Strategy

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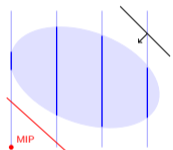
- **Goal:** bring together the primal and dual side, i.e. find the optimal solution and prove that it is optimal
 - Sometimes finding a solution within a certain proven gap is enough
- A brief overview of algorithms for convex MINLPs
- Spatial branch and bound for nonconvex MINLPs

Algorithms for Convex MINLP: Overview



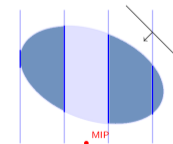
Outer Approximation:

- Solve **MILP relaxations** and **NLP subproblems**
- Add gradient cuts at solutions of NLP subproblems
- Uses the equivalence of MINLP to MILP



Extended Cutting Planes:

- Solve MILP relaxations
- Add gradient cuts at solutions of **MILP relaxations**



Branch and Bound:

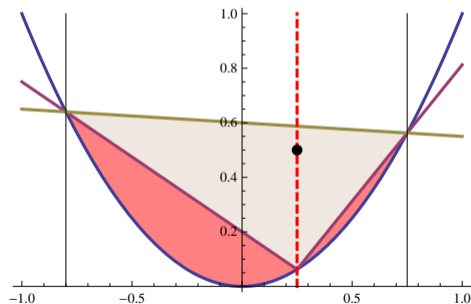
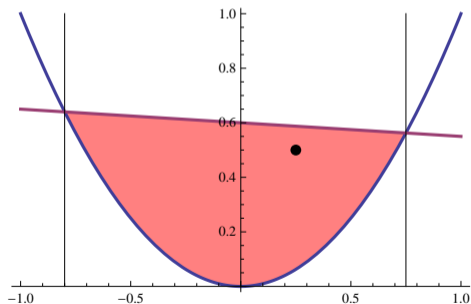
- Generalization of MILP B&B
- The continuous relaxation is **nonlinear** (but convex)
- Different choices between LP and NLP relaxations

Algorithms for Nonconvex MINLP: Spatial Branching

- Recall: variable bounds determine the convex relaxation, e.g.,

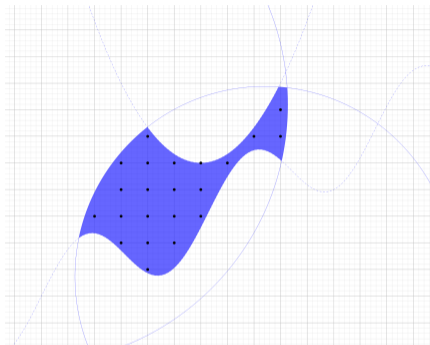
$$x^2 \leq \ell^2 + \frac{u^2 - \ell^2}{u - \ell}(x - \ell) \quad \forall x \in [\ell, u]$$

- Branch on variables in **violated nonconvex** constraints to improve relaxations



Spatial Branch and Bound for MINLP

- Solve a **relaxation** \rightarrow lower bound
- **Branch** on a suitable variable (integer, or continuous in a violated nonconvex constraint)
- If a solution is **integer feasible** and **satisfies nonlinear constraints** \rightarrow upper bound
- **Discard** parts of the tree that are infeasible or where lower bound \geq best known upper bound
- **Repeat** until gap is below given tolerance

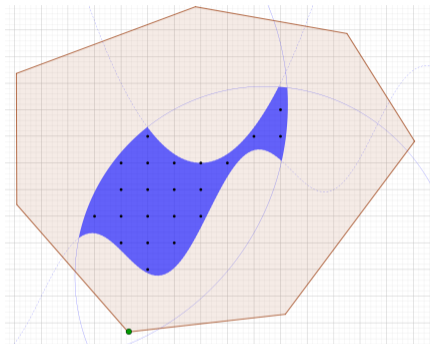


Smaller domains \rightarrow improved relaxations \rightarrow improved bounds.

In practice algorithms also include presolving, heuristics, propagation, etc.

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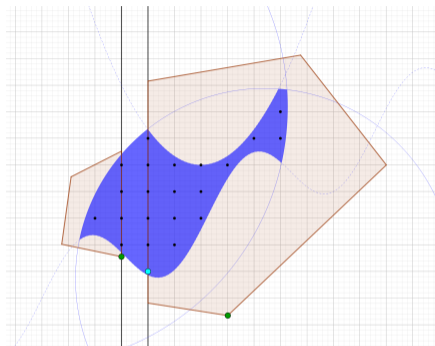


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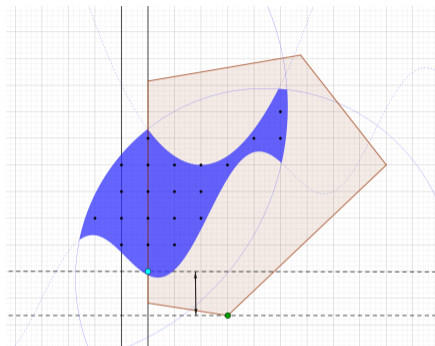


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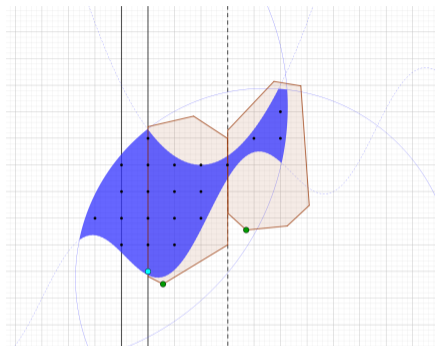


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- **Repeat** until gap is below given tolerance



Smaller domains \rightarrow improved relaxations \rightarrow improved bounds.

In practice algorithms also include presolving, heuristics, propagation, etc.

MINLP in SCIP

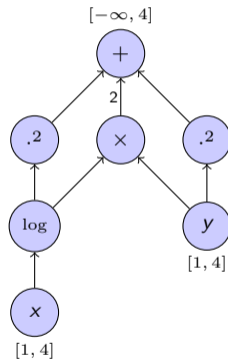
MINLP in SCIP

- SCIP is a constraint integer programming (CIP) solver
- CIP is a generalization of MILP allowing for arbitrary constraints as long as a tractable relaxation can be built
- CIP includes MINLP (and a few other problem classes)

- SCIP implements **LP-based spatial B&B**
- Convex relaxations are constructed via the **auxiliary variable method**
- The handling of nonlinear constraints is based on **expression graphs**
- The nonlinear constraint handler coordinates sub-plugins handling various nonlinearities

Expression Trees in SCIP

$$\log(x)^2 + 2 \log(x)y + y^2$$

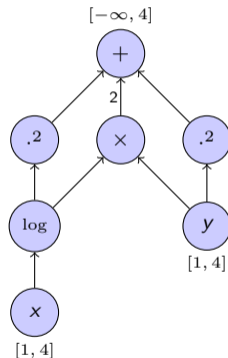


Expression Trees in SCIP

Operators:

- variable index, constant
- $+$, $-$, $*$, \div
- \cdot^2 , $\sqrt{\cdot}$, \cdot^p ($p \in \mathbb{R}$), \cdot^n ($n \in \mathbb{Z}$), $x \mapsto x|x|^{p-1}$ ($p > 1$)
- exp, log
- min, max, abs
- \sum , \prod , affine-linear, quadratic, signomial
- (user)

$$\log(x)^2 + 2 \log(x)y + y^2$$

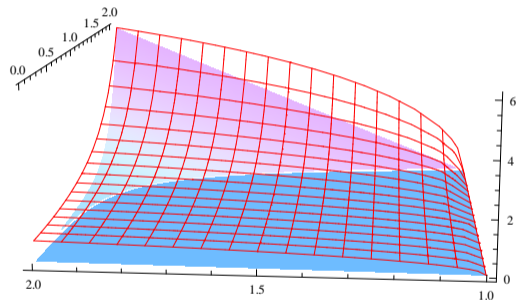


Reformulation (During Presolve)

Goal: **reformulate constraints** such that only **elementary cases** (convex, concave, odd power, quadratic) remain. Implements the **auxiliary variable method**.

Example:

$$g(x) = \sqrt{\exp(x_1^2) \ln(x_2)}$$



Introduces **new variables and new constraints**.

Reformulation (During Presolve)

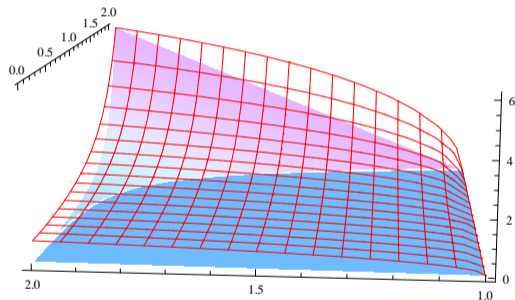
Goal: **reformulate constraints** such that only **elementary cases** (convex, concave, odd power, quadratic) remain. Implements the **auxiliary variable method**.

Example:

$$g(x) = \sqrt{\exp(x_1^2) \ln(x_2)}$$

Reformulation:

$$\begin{aligned}g &= \sqrt{y_1} \\ y_1 &= y_2 y_3 \\ y_2 &= \exp(y_4) \\ y_3 &= \ln(x_2) \\ y_4 &= x_1^2\end{aligned}$$



Introduces **new variables and new constraints**.

Problem with classic approach

Explicit reformulation of constraints ...

- ... **loses the connection to the original problem.**
- ... **loses distinction between original and auxiliary variables.** Thus, we may branch on auxiliary variables.
- ... **prevents simultaneous exploitation of overlapping structures.**

SCIP's Handling of Reformulations

- Avoid explicit split-up of constraints
- Introduce **extended formulation** as **annotation to the original** formulation
- Use extended formulation for **relaxation**
- Use original formulation for **feasibility checking**
- To resolve infeasibility in original constraints, tighten relaxation of extended formulation
- The original formulation is kept
- This avoids wrong feasibility checks

Practical Topics

Impact of Modeling

The choice of formulation makes a difference.

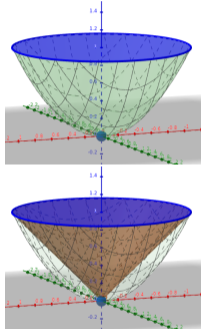
Example: x and y contained in circle of radius c if $z = 1$ and are both zero if $z = 0$.

One could model this as:

$$\begin{aligned}x^2 + y^2 &\leq cz \\ x, y &\in \mathbb{R}, z \in \{0, 1\}\end{aligned}$$

Or as:

$$\begin{aligned}x^2 + y^2 &\leq cz^2 \\ x, y &\in \mathbb{R}, z \in \{0, 1\}\end{aligned}$$



These **describe the same feasible set** ($z^2 = z$ if $z \in \{0, 1\}$). But the **second** formulation leads to a **tighter continuous relaxation** ($z^2 < z$ if $z \in (0, 1)$).

How to Experiment

- Performance variability
 - **Significant changes in performance** caused by **small changes in model/algorithm**
 - Occurs in MILP, but tends to be even more pronounced in MINLP
- Obtaining more reliable results
 - If possible and makes sense, use **large** and **heterogeneous** testsets
 - Take advantage of performance variability: model permutations (reordering variables and constraints) can help against random effects (for example, in SCIP this is controlled by a parameter)
- Using solver statistics
 - Information on tree **nodes**, primal and dual **bounds**, effects of solver **components**
 - Helpful for finding bottlenecks
- Isolating feature effects
 - **Turn off some components** to get rid of some random effects...
 - or to analyse interaction: some component might make the feature redundant, etc.

Recap

- MINLPs combine integrality and nonlinearity
- Algorithms are based on finding and improving primal and dual bounds
- Primal bounds are found by heuristics; there are many extensions of MILP techniques
- Dual bounds are found via relaxations (usually convex or linear)
- Spatial branch and bound solves nonconvex MINLPs globally by also branching on continuous variables

Questions?