

GAMS

Stefan Vigerske
stefan@gams.com



October 2nd, 2015, CO@Work, Berlin

Material to this lecture: <http://co-at-work.zib.de/files/gams/>

- ▶ **CO@Work virtual machines:** GAMS is installed (run gams)
- ▶ **Download** GAMS system: <http://www.gams.com/download>
- ▶ **Evaluation license** (valid until 7.1.2016):
<http://co-at-work.zib.de/files/gams/gamslice.txt>
(already installed on 32bit VM)

Material to this lecture: <http://co-at-work.zib.de/files/gams/>

- ▶ **CO@Work virtual machines:** GAMS is installed (run gams)
- ▶ **Download** GAMS system: <http://www.gams.com/download>
- ▶ **Evaluation license** (valid until 7.1.2016):
<http://co-at-work.zib.de/files/gams/gamslice.txt>
(already installed on 32bit VM)

Next 85 minutes:

- ▶ A short overview on GAMS.
- ▶ ~~GAMS syntax and programming.~~ (gamsbasics.pdf)
- ▶ Equitable coach problem in GAMS.
- ▶ Exercise: Run and modify equitable coach problem.
- ▶ “Hiding GAMS”: Calling GAMS from Python.
- ▶ Exercise: Cutting Stock by Column Generation in Python.

- ▶ **Roots:** World Bank, 1976
- ▶ Went **commercial** in 1987
- ▶ **Locations**
 - ▶ GAMS Development Corporation (Washington, D.C.)
 - ▶ GAMS Software GmbH (Germany)



Washington



Braunschweig

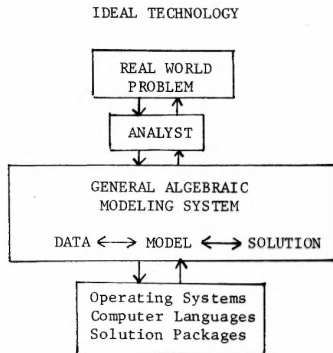


Frechen (Cologne)



ZIB, Berlin

- ▶ \approx **16 employees** (7 development, 5 sales, ...)
- ▶ **Product:** The **G**eneral **A**lgebraic **M**odeling **S**ystem



- RESULT:
- Limited drain of resources
 - Same representation of models for humans and machines
 - Model representation is also model documentation



- 1976 GAMS idea is presented at ISMP, Budapest
- 1978 Phase I: GAMS supports **linear programming**. Supports **Mainframes** and **Unix** Workstations
- 1979 Phase II: GAMS supports **nonlinear programming**
- 1987 GAMS becomes a commercial product
- 1988 First **PC System** (16 bit)
- 1988 Alex Meeraus, the initiator of GAMS and founder of GAMS Development Corporation, is awarded INFORMS Computing Society Prize
- 1990 **32 bit Dos Extender**
- 1990 GAMS moves to Georgetown, Washington, D.C.
- 1991 **Mixed Integer Non-Linear** Programs capability (DICOPT)
- 1994 GAMS supports **mixed complementarity** problems
- 1995 MPSGE language is added for CGE modeling
- 1996 European branch opens in Germany
- 1998 **32 bit native Windows**
- 1998 **Stochastic programming** capability (OSL/SE, DECIS)
- 1999 Introduction of the GAMS **Integrated development environment** (IDE)
- 2000 GAMS World initiative started
- 2001 **GAMS Data Exchange** (GDX) is introduced
- 2002 GAMS is listed in OR/MS 50th Anniversary list of milestones
- 2003 **Conic programming** is added
- 2003 **Global optimization** in GAMS
- 2004 Support for Quadratic Constrained programs
- 2005 Support for **64 bit PC** Operating systems
- 2006 GAMS supports parallel **grid computing**
- 2007 GAMS supports open-source solvers from COIN-OR
- 2008 Support for 32 and 64 bit **Mac OS X**
- 2009 GAMS supports extended mathematical programs (EMP)
- 2010 GAMS is awarded the company award of the German Society of Operations Research (GOR)
- 2012 The Winners of the 2012 INFORMS Impact Prize included Alexander Meeraus. The prize was awarded to the originators of the five most important algebraic modeling languages.

- ▶ more than **11500 licenses**
- ▶ **50% academic** users, **50% commercial**

Application Areas:

- | | |
|---------------------------|-------------------------------|
| ▶ Agricultural Economics | ▶ Applied General Equilibrium |
| ▶ Chemical Engineering | ▶ Economic Development |
| ▶ Econometrics | ▶ Energy |
| ▶ Environmental Economics | ▶ Engineering |
| ▶ Finance | ▶ Forestry |
| ▶ International Trade | ▶ Logistics |
| ▶ Macro Economics | ▶ Military |
| ▶ Management Science/OR | ▶ Mathematics |
| ▶ ... | |

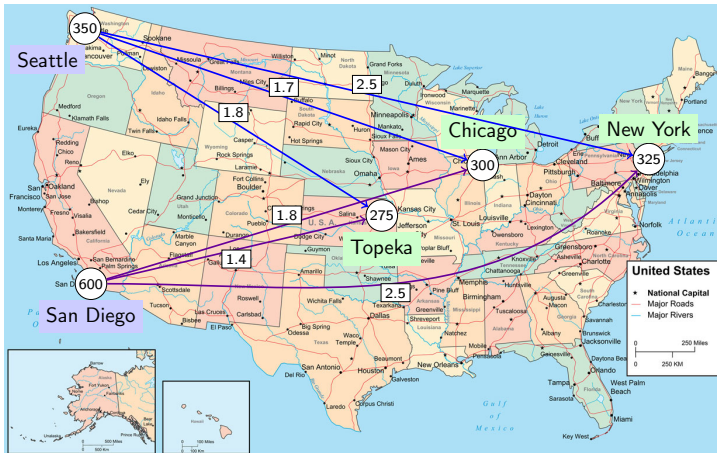
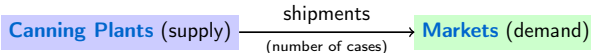


Declarative Language:

- ▶ Similar to **mathematical notation**
- ▶ Few **basic language elements**: sets, parameters, variables, equations, models
- ▶ Model is executable (algebraic) description of the problem

Imperative Elements:

- ▶ **Control flow** statements: loops, for, if, ...
- ▶ build algorithms within GAMS
- ▶ **exchange data** with other systems



Minimize Transportation cost
 subject to Demand satisfaction at markets
 Supply constraints

i canning plants
 j markets
 a_i capacity of plants
 b_j demand at markets
 $c_{i,j}$ transportation cost per case
 $x_{i,j}$ cases to ship from i to j

$$\min \sum_{i,j} c_{i,j} x_{i,j}$$

$$\text{s.t.} \quad \sum_j x_{i,j} \leq a_i \quad \forall i$$

$$\sum_i x_{i,j} \geq b_j \quad \forall j$$

$$x_{i,j} \geq 0 \quad \forall i,j$$

Sets

```
i  "canning plants",
j  "markets";
```

Parameters

```
a(i)  "capacity of plant i in cases",
b(j)  "demand at market j in cases",
c(i,j) "transport cost in 1000$ per case";
```

Variables

```
x(i,j) "shipment quantities in cases",
z      "total transportation costs in 1000$"
```

Positive Variable x;

Equations

```
cost      "define objective function",
supply(i) "observe supply limit at plant i",
demand(j) "satisfy demand at market j";
```

```
cost ..      z =e= sum((i,j), c(i,j)*x(i,j));
supply(i) .. sum(j, x(i,j)) =l= a(i);
demand(j) .. sum(i, x(i,j)) =g= b(j);
```

```
Model transport / all / ;
```



GAMS is not a Solver!

GAMS: Model building and interaction with solvers and environment.

Solver: Solve an instance (instantiation of a model with data) using mathematical optimization.

- ▶ Major commercial and academic **solvers integrated**:
31 solvers, half of them actively developed/updated
- ▶ Average number of commercial **solvers per license**:
 - ▶ Academic clients: 2.9
 - ▶ Commercial clients: 2.2
- ▶ **Switch between solvers** with one statement:
`option solver = scip;` (since GAMS 24.5)

	LP	MIP	NLP	MCP	MPEC	CNS	DNLP	MINLP	QCP	MIQCP	Stoch.	Global
ALPHAECIP								x		x		
ANTIGONE 1.1			x			x	x	x	x	x		x
BARON 15.8	x	x	x			x	x	x	x	x		x
BDMLP	x	x										
BONMIN 1.8								x		x		
CBC 2.9	x	x										
CONOPT 3	x		x			x	x		x			
COUENNE 0.5			x			x	x	x	x	x		x
CPLEX 12.6	x	x							x	x		
DECIS	x										x	
DICOPT								x		x		
GUROBI 6.0	x	x							x	x		
IPOPT 3.12	x		x			x	x		x			
KNITRO 9.1	x		x			x	x	x	x	x		
LGO	x		x				x		x			(x)
LINDO 9.0	x	x	x				x	x	x	x	x	x
LOCALSOLVER 5.5	x	x	x			x	x	x	x	x		
MILES				x								
MINOS	x		x			x	x		x			
MOSEK 7	x	x	x				x		x	x		
MSNLP			x				x		x			(x)
NLPEC				x	x							
OQNLP			x				x	x	x	x		(x)
PATH				x		x						
SBB								x		x		
SCIP 3.2		x	x			x	x	x	x	x		x
SNOPT	x		x			x	x		x			
SOPLEX 2.2	x											
SULUM 4.3	x	x										
XA	x	x										
XPRESS 28.01	x	x							x	x		

Independence of Model and Platform

GAMS

Supported Platforms:



Solver/Platform availability - 24.5							
	x86 32bit	x86 64bit	x86 64bit	x86 64bit	x86 64bit	Sparc 64bit	IBM Power 64bit
	MS Windows	MS Windows	Linux	MacOS X	SOLARIS	SOLARIS	AIX
ALPHAEC	✓	✓	✓	✓	✓	✓	✓
ANTIGONE 1.1	✓	✓	✓	✓			
BARON 15.8	✓	✓	✓	✓			
BOMLP	✓	✓	✓	✓	✓		✓
BONMIN 1.8	✓	✓	✓	✓	✓		
CBC 2.9	✓	✓	✓	✓	✓		
CONOPT 3	✓	✓	✓	✓	✓		✓
COUENNE 0.5	✓	✓	✓	✓	✓		
CPLEX 12.6	✓	✓	✓	✓	✓		✓
DÉCIS	✓	✓	✓	✓	✓		
DICOPT	✓	✓	✓	✓	✓		✓
GLMGO 2.3	✓	✓	✓	✓	✓		
GUROBI 6.0	✓	✓	✓	✓	✓		✓
GUSS	✓	✓	✓	✓	✓		✓
IPOPT 3.12	✓	✓	✓	✓	✓		✓
KESTREL	✓	✓	✓	✓	✓		✓
KNITRO 9.1	✓	✓	✓	✓	✓		
LGO	✓	✓	✓	✓	✓		
LINDO 9.0	✓	✓	✓	✓	✓		
LINDO GLOBAL 9.0	✓	✓	✓	✓	✓		
LOCALSOLVER 5.5	✓	✓	✓	✓	✓		
MILES	✓	✓	✓	✓	✓		✓
MILVOS	✓	✓	✓	✓	✓		
MOSEK 7	✓	✓	✓	✓	✓		
MSNLPL	✓	✓	✓	✓	✓		
NLPEC	✓	✓	✓	✓	✓		✓
OQNLP	✓	32bit	✓	✓	✓		
PATH	✓	✓	✓	✓	✓		✓
SBB	✓	✓	✓	✓	✓		✓
SCIP 3.2	✓	✓	✓	✓	✓		
SNOPT	✓	✓	✓	✓	✓		✓
SOLPLEX 2.2	✓	✓	✓	✓	✓		
SULLM 4.3	✓	✓	✓	✓	✓		
XA	✓	✓	✓	✓	✓		
XPRESS 28.01	✓	✓	✓	✓	✓		✓

Tools/Platform availability - 24.5							
	x86 32bit	x86 64bit	x86 64bit	x86 64bit	x86 64bit	Sparc 64bit	IBM Power 64bit
	MS Windows	MS Windows	Linux	MacOS X	SOLARIS	SOLARIS	AIX
ASK	✓	32bit	✓	✓	✓	✓	✓
BIB2GMS	✓	✓	✓	✓	✓	✓	✓
CHK4UPD	✓	✓	✓	✓	✓	✓	✓
CHOLESKY	✓	✓	✓	✓	✓	✓	✓
CSDP	✓	✓	✓	✓	✓	✓	✓
CSV2GDX	✓	✓	✓	✓	✓	✓	✓
EIGENVALUE	✓	✓	✓	✓	✓	✓	✓
EIGENVECTOR	✓	✓	✓	✓	✓	✓	✓
ENDECRYPT	✓	✓	✓	✓	✓	✓	✓
GAMSICE	✓	32bit	✓	✓	✓	✓	✓
GAMS POSIX Utilities ¹¹	✓	✓	✓	✓	✓	✓	✓
GDX2ACCESS	✓	32bit	✓	✓	✓	✓	✓
GDX2HAR	✓	32bit	✓	✓	✓	✓	✓
GDX2SQLITE	✓	✓	✓	✓	✓	✓	✓
GDX2VEDA	✓	✓	✓	✓	✓	✓	✓
GDX2XLS	✓	32bit	✓	✓	✓	✓	✓
GDXCOPY	✓	✓	✓	✓	✓	✓	✓
GDXDIFF	✓	✓	✓	✓	✓	✓	✓
GDXDUMP	✓	✓	✓	✓	✓	✓	✓
GDXMERGE	✓	✓	✓	✓	✓	✓	✓
GDXMRW	✓	✓	✓	✓	✓	✓	✓
GDXRANK	✓	✓	✓	✓	✓	✓	✓
GDXRENAME	✓	✓	✓	✓	✓	✓	✓
GDXRRW	✓	✓	✓	✓	src only	src only	src only
GDXTROLL	✓	✓	✓	✓	✓	✓	✓
GDXVIEWER	✓	32bit	✓	✓	✓	✓	✓
GDXXRW	✓	32bit	✓	✓	✓	✓	✓
GMSUNZIP	✓	✓	✓	✓	✓	✓	✓
HAR2GDX	✓	32bit	✓	✓	✓	✓	✓
IDECMS	✓	32bit	✓	✓	✓	✓	✓
INVERT	✓	✓	✓	✓	✓	✓	✓
MCFILTER	✓	✓	✓	✓	✓	✓	✓
MDB2GMS	✓	32bit	✓	✓	✓	✓	✓
MODEL2TEX	✓	✓	✓	✓	✓	✓	✓
MP52GMS	✓	✓	✓	✓	✓	✓	✓
MSAPPA/AIL	✓	32bit	✓	✓	✓	✓	✓
SCENRED	✓	✓	✓	✓	✓	✓	✓
SCENRED2	✓	✓	✓	✓	✓	✓	✓
SHELLEXECUTE	✓	32bit	✓	✓	✓	✓	✓
SOLGMS	✓	32bit	✓	✓	✓	✓	✓

Online: <http://www.gams.com/help> (with search)

Offline: <GAMS system directory>/docs/index.html (no search, use grep!)

- ▶ **GAMS – A User's Guide:** Tutorial, Basics, Advanced Topics
- ▶ **McCarl (Expanded) GAMS User Guide**
- ▶ **Solver Manuals**
- ▶ **Tools Manuals**
- ▶ **APIs:** Tutorials and Reference Manuals
- ▶ **Release Notes**

Tutorial Videos: <http://www.youtube.com/user/GAMSLessons>

Support wiki: <http://support.gams.com/doku.php>

Discussion group: <http://www.gamsworld.org/>

Online: <http://www.gams.com/modlibs>

Offline: gamslib, apilib, datalib, emplib, testlib tools

► GAMS Model Library

- representing interesting and sometimes classic problems
- illustrating GAMS modeling capabilities

Online: <http://www.gams.com/modlibs>

Offline: gamslib, apilib, datalib, emplib, testlib tools

- ▶ **GAMS Model Library**

- ▶ representing interesting and sometimes classic problems
- ▶ illustrating GAMS modeling capabilities

- ▶ **GAMS API Library**

- ▶ scripts to compile and execute GAMS API examples

- ▶ **GAMS Data Utilities Library**

- ▶ demonstrate utilities to interface GAMS with other applications

- ▶ **GAMS EMP Library**

- ▶ illustrate and test capabilities of extended mathematical programming facility

- ▶ **Contributed Libraries:**

- ▶ FINLIB – financial optimization models (by Consiglio, Nielsen and Zenios)
- ▶ NOALIB – nonlinear optimization applications models (by Neculai Andrei)

- ▶ **GAMS Testlib Library**

- ▶ testing and quality control

The Equitable Coach Problem



What is important?

- ▶ How many players are in the team? How many on the field?
- ▶ What is the formation that the team will play? 2-3-1, 2-2-2 or 3-1-2?
- ▶ How many interchanges do you want to perform?
- ▶ What are the positional preferences of the players?
- ▶ Are you equitable for only a single game or across the two matches?
- ▶ **Complicated:** Does winning the first match impact your decision about being equitable for the second?

Burt, Maher, Witzig – Modelling

11 / 102



Problem definition



OBJECTIVE: For a single game identify the equitable number of playing minutes for each player given that there is only one goal keeper and **players have positional preferences.**

- ▶ 7 players on the field, 12 players in a team
- ▶ 20 minutes games \Rightarrow 140 player minutes
- ▶ 4 different position types, Goal Keeper, Defence, Midfield and Striker.

Solution is not as easy. Need a more sophisticated mathematical model.

Burt, Maher, Witzig – Modelling

12 / 102

The Equitable Coach Problem: disaggregated model



We now consider the version of the ECP where we allocate players to positions at particular times.

- $pr_{i,t}$ [binary] is 1 if player r plays in position i in play period t ;
- e_r^+ [continuous] is the number of periods player r plays above the average;
- e_r^- [continuous] is the number of periods player r plays below the average.
- R is the set of players;
- N is the set of positions;
- T is the total number of game minutes.

Burt, Maher, Witzig – Modelling

18 / 102



The Equitable Coach Problem: disaggregated model



$$\begin{aligned}
 ECPd: \quad & \min \sum_r (e_r^+ + e_r^-) \\
 & \sum_r pr_{r,i,t} = 1 \quad \forall t \in T, i \in N, \quad (8) \\
 & \sum_i pr_{r,i,t} \leq 1 \quad \forall t \in T, r \in R, \quad (9) \\
 & \sum_{t,i} pr_{r,i,t} - \sum_{t,i} \frac{pr_{r,i,t}}{|R|} = e_r^+ - e_r^- \quad \forall r \in R, \quad (10) \\
 & e_r^+ \in [0, 1] \quad \forall r \in R, i \in N, t \in T, \\
 & e_r^+, e_r^- \in \mathbb{R}_{\geq 0}.
 \end{aligned}$$

Burt, Maher, Witzig – Modelling

19 / 102

- ▶ 7 players on the field, 12 players in a team
- ▶ 20 minutes games \Rightarrow 140 player minutes
- ▶ 4 different position types, Goal Keeper, Defence, Midfield and Striker.

OBJECTIVE: For a single game identify the **positions of players at particular times** given that players have positional preferences and such that **playing times are as equitable as possible**.

- ▶ 7 players on the field, 12 players in a team
- ▶ 20 minutes games \Rightarrow 140 player minutes
- ▶ 4 different position types, Goal Keeper, Defence, Midfield and Striker.

OBJECTIVE: For a single game identify the **positions of players at particular times** given that players have positional preferences and such that **playing times are as equitable as possible**.

$$\begin{aligned}
 & x_{r,i,t} \text{ is 1 if player } r \text{ plays in position } i \text{ in play period } t; \\
 & \epsilon_r^{+/-} \text{ is the number of periods player } r \text{ plays above/below the average;} \\
 & R \text{ is the set of players;} \\
 & R_i \text{ is the set of players than can play on position } i; \\
 & P \text{ is the set of positions;} \\
 & T \text{ is the total number of game minutes.}
 \end{aligned}
 \quad
 \begin{aligned}
 & \min \sum_r (\epsilon_r^+ + \epsilon_r^-) \\
 & \sum_{r \in R_i} x_{r,i,t} = 1 \quad \forall t \in T, i \in P, \\
 & \sum_{i: r \in R_i} x_{r,i,t} \leq 1 \quad \forall t \in T, r \in R, \\
 & \sum_{t, i: r \in R_i} x_{r,i,t} - \sum_{t, i, r' \in R_i} \frac{x_{r',i,t}}{|R|} = \epsilon_r^+ - \epsilon_r^- \quad \forall r \in R, \\
 & x_{r,i,t} \in [0, 1] \quad \forall r \in R, i \in N, t \in T, \\
 & \epsilon_r^+, \epsilon_r^- \in \mathbb{R}_{\geq 0}.
 \end{aligned}$$

```
param gametime      := 20;
param nsubperiods   := 6;

# formation
param nkeeper       := 1;
param ndefender     := 2;
param nmidfielder   := 3;
param nstriker      := 1;

# set number of positions
param npositions    := nkeeper + ndefender + nmidfielder + nstriker;

# sets of players, positions, and minutes per game
set Player          := {
    "Josh", "Simon", "Jordy", "Chris", "Andy", "Richie",
    "Tritto", "Guy", "Neil", "Justin", "Steve", "Cory"
};
set Classes         := { "K", "D", "M", "S" };
set Positions       := { 1 .. npositions };
set Periods         := { 1 .. nsubperiods };

...

var x[Player * Positions * Periods] binary;
var eplus[Player] real;
var eminus[Player] real;

# the model:
# 1.) objective function
# 2.) ensure that each position is used at each point in time
# 3.) ensure that each player plays on at most one position per time
# 4.) calculate the minutes over/below the average
minimize cost: sum <r> in Player : gametime/nsubperiods * (eplus[r] + eminus[r]);

subto c1: forall <i> in Positions :
    forall <t> in Periods : sum <r> in Player with pos[r,GetClass(i)] == 1: x[r,i,t] == 1;

subto c2: forall <r> in Player :
    forall <t> in Periods : sum <i> in Positions with pos[r,GetClass(i)] == 1: x[r,i,t] <= 1;
```

\$title Equitable Coach Problem

** timing parameters*

** param gametime := 20;*

scalar gametime / 20 /;

** param nsubperiods := 6;*

\$set nsubperiods 6

** formation 1-2-3-1*

** param nkeeper := 1;*

** param ndefender := 2;*

** param nmidfielder := 3;*

** param nstriker := 1;*

scalar nkeeper / 1 /;

scalar ndefender / 2 /;

scalar nmidfielder / 3 /;

scalar nstriker / 1 /;

** set number of positions*

** param npositions := nkeeper + ndefender + nmidfielder + nstriker;*

\$set npositions 7

```
* sets of players, positions, and minutes per game
*set Player          := {
*   "Josh",    "Simon", "Jordy", "Chris", "Andy", "Richie",
*   "Tritto",  "Guy",    "Neil",  "Justin", "Steve", "Cory"
*};
set r "Player" /
    "Josh",    "Simon", "Jordy", "Chris", "Andy", "Richie",
    "Tritto",  "Guy",    "Neil",  "Justin", "Steve", "Cory"
/;

*set Classes        := { "K", "D", "M", "S" };
*set Positions      := { 1 .. npositions };
*set Periods        := { 1 .. nsubperiods };
set c "Position classes" / K, D, M, S /;
set p "Positions" / 1 * %npositions% /;
set t "Periods" / 1 * %nsubperiods% /;
```

```
* entry (i,j) is 1 <=> player i can play on position j
```

```
$ontext
```

```
param pos[Player * Classes] := | "K", "D", "M", "S" |
                                | "Josh"  | 0, 0, 1, 1 |
                                | "Simon" | 0, 0, 0, 1 |
```

```
...
```

```
                                | "Steve" | 0, 1, 0, 0 |
                                | "Cory"  | 1, 0, 0, 0 |;
```

```
$offtext
```

```
Table pos(r,c)
```

	K	D	M	S
Josh			1	1
Simon				1
Jordy				1
Chris	1	1		
Andy		1		
Richie	1	1		
Tritto	1			
Guy	1			
Neil		1		
Justin		1		
Steve	1			
Cory	1			

```
;
```

\$ontext

```
defstrg GetClass(i) := if i <= nkeeper then "K"  
                    else  
                        if i <= nkeeper + ndefender then "D"  
                        else  
                            if i <= nkeeper + ndefender + nmidfielder  
                                else  
                                    "S"  
                                end  
                            end  
                        end  
                    end;  
end;
```

\$offtext

```
set pos2class(p,c);  
loop(p,  
    if( p.val <= nkeeper,  
        pos2class(p,'K') = yes;  
    else if( p.val <= nkeeper + ndefender,  
        pos2class(p,'D') = yes;  
    else if( p.val <= nkeeper + ndefender + nmidfielder,  
        pos2class(p,'M') = yes;  
    else  
        pos2class(p,'S') = yes;  
    )))  
);
```



```
set canplay(r,p) "whether player r can play in position p";  
* r can play at p if pos(r,c) is 1 for the position class that p belongs to  
canplay(r,p) = (sum(c$pos2class(p,c), pos(r,c)) = 1);
```

```
display pos, pos2class, canplay;
```

```
* variable definition  
*var x[Player * Positions * Periods] binary;  
*var eplus[Player] real;  
*var eminus[Player] real;  
Binary Variable x(r,p,t);  
Positive Variables eplus(r), eminus(r);
```

```
* the model:  
* 1.) objective function  
* 2.) ensure that each position is used at each point in time  
* 3.) ensure that each player plays on at most one position per time  
* 4.) calculate the minutes over/below the average
```

```
Variable inquiry;
```

```
Equations obj, c1, c2, c3;
```

```

* minimize cost: sum <r> in Player : gametime/nsubperiods * (eplus[r]+eminus[r]);
obj.. inquiry =E= sum(r, gametime/%nsubperiods% * (eplus(r) + eminus(r)));

*subto c1: forall <i> in Positions : forall <t> in Periods :
*      sum <r> in Player with pos[r,GetClass(i)] == 1: x[r,i,t] == 1;
c1(p,t).. sum(r$canplay(r,p), x(r,p,t)) =E= 1;

*subto c2: forall <r> in Player : forall <t> in Periods :
*      sum <i> in Positions with pos[r,GetClass(i)] == 1: x[r,i,t] <= 1;
c2(r,t).. sum(canplay(r,p), x(r,p,t)) =L= 1;

*subto c3: forall <r> in Player :
*      sum <i,t> in Positions * Periods with pos[r,GetClass(i)] == 1 : x[r,i,t]
* -sum <q,i,t> in Player * Positions * Periods with
*      pos[q,GetClass(i)] == 1 : (x[q,i,t]/card(Player)) == eplus[r]-eminus[r]
alias(r,q);
c3(r).. sum((p,t)$canplay(r,p), x(r,p,t))
      - sum((q,p,t)$canplay(q,p), x(q,p,t)/card(r)) =E= eplus(r)-eminus(r)

* declare model ECPd, consisting of all equations seen so far
Model ECPd / all /;

* solve ECPd (minimizing inquiry) as MIP to optimality (gap tolerance 0)
option optcr = 0;

```

Solve ECPd minimizing inquiry using MIP;

TASK: Extend the equitable coach problem model by restricting the number of position changes for all players to 3.

- ▶ Get http://co-at-work.zib.de/files/gams/ECP_exercise.gms
- ▶ Open the file in your favorite text editor and read the more detailed instructions at the beginning.
- ▶ Run the model: `gams ECP_exercise.gms errmsg=1`
- ▶ In case of compilation or execution errors, e.g.,

```
--- Starting compilation
--- ECP_exercise.gms(111) 3 Mb 3 Errors
*** Status: Compilation error(s)
--- Job ECP_exercise.gms Stop 10/01/15 20:15:47 elapsed 0:00
```

see listing file (ECP_exercise.lst):

```
90  c2(r,t).. sum(canplay(r,p), x(r,p,t)) <= 1;
****                                     $37
**** 37  '=l=' or '=e=' or '=g=' operator expected
```

GD_X – GAMS Data Exchange:

- ▶ **Binary data** format for fast exchange of data with GAMS
- ▶ Stores **sets**, **parameters**, **values of variables/equations** with domain information, but no symbolic information (e.g., equation algebra)
- ▶ **Consistency**: no duplicates or contradictions
- ▶ **Platform independent**
- ▶ Can be **compressed**

GDX – GAMS Data Exchange:

- ▶ **Binary data** format for fast exchange of data with GAMS
- ▶ Stores **sets**, **parameters**, **values of variables/equations** with domain information, but no symbolic information (e.g., equation algebra)
- ▶ **Consistency**: no duplicates or contradictions
- ▶ **Platform independent**
- ▶ Can be **compressed**
- ▶ **Read and write** in GAMS:
 - ▶ on command line: `parameter gdx=<filename>`
 - ▶ during compilation: `$gdxin`, `$gdxout`, `$load`, `$unload`, ...
 - ▶ during execution: `execute_load`, `execute_unload`, ...
- ▶ **Tools and APIs** to read and write from other environments:
 - ▶ `gdxdump`, `csv2gdx`, `gdxviewer` (win only), ...
 - ▶ Matlab, MS Access, MS Excel, ODBC/SQL, R, SQLite
 - ▶ low-level APIs for C, C++, C#, Delphi, Fortran, Java, Python, VBA, VB.NET
 - ▶ high-level APIs for C#, Java, and Python

Creating input:

- ▶ Read GAMS model from file or string (APIs have no modeling capability!)
- ▶ Create input data (wrapper around GDX API)

Callout to GAMS:

- ▶ Set GAMS options (e.g., solver to use, gap limit, ...)
- ▶ Start GAMS run and store results in GDX file

Processing results:

- ▶ Read and process result data (wrapper around GDX API)

Creating input:

- ▶ Read GAMS model from file or string (APIs have no modeling capability!)
- ▶ Create input data (wrapper around GDX API)

Callout to GAMS:

- ▶ Set GAMS options (e.g., solver to use, gap limit, ...)
- ▶ Start GAMS run and store results in GDX file

Processing results:

- ▶ Read and process result data (wrapper around GDX API)

Main Classes:

- ▶ GAMSWorkspace: Environment
- ▶ GAMSJob: Storing and running a model
- ▶ GAMSDatabase: In-memory representation of data ("GDX in-memory")
- ▶ GAMSOptions: Manipulation of GAMS options
- ▶ GAMSModelInstance: Hot-start solve for a sequence of closely related problems (changes in numerical values of matrix, objective, sides, bounds only)

<GAMS system directory>/apifiles/Python/transport1.py:

```
from gams import *
import sys

if len(sys.argv) > 1:
    ws = GamsWorkspace(system_directory = sys.argv[1])
else:
    ws = GamsWorkspace()

ws.gamslib("trnsport")      # get 'trnsport' from GAMS model library
t1 = ws.add_job_from_file("trnsport.gms")

opt = ws.add_options()
opt.all_model_types = "soplex"
t1.run(opt)

for rec in t1.out_db["x"]:
    print("x(" + rec.key(0) + "," + rec.key(1) + "): " + \
          "level=" + str(rec.level) + " marginal=" + str(rec.marginal))

x(seattle,new-york): level=50.0 marginal=0.0
x(seattle,chicago): level=300.0 marginal=0.0
x(seattle,topeka): level=0.0 marginal=0.036000000000000004
x(san-diego,new-york): level=275.0 marginal=0.0
x(san-diego,chicago): level=0.0 marginal=0.00900000000000000008
x(san-diego,topeka): level=275.0 marginal=0.0
```


Problem:

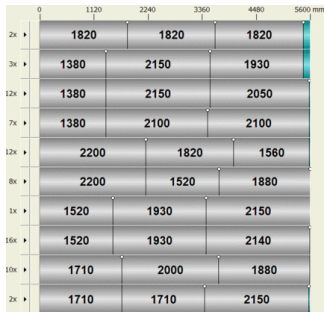
[Gilmore, Gomory, 1961, 1963]

- ▶ Cut out some paper products of different sizes from large raw paper rolls, in order to meet a customer's order.
- ▶ The objective is to minimize the required number of the paper rolls.

Example:

- ▶ Raw paper roll width: 5600mm
- ▶ Customer's Order:

Width	Rolls	Width	Rolls
1380	22	2000	10
1520	25	2050	12
1560	12	2100	14
1710	14	2140	16
1820	18	2150	18
1880	18	2200	20
1930	20		



Source: Wikipedia

Width of raw paper roll: W

$$W = 5600$$

Customer order:

- ▶ Set I of widths.
- ▶ Widths $w_i \in \mathbb{N}$, $i \in I$.
- ▶ Required number of rolls for each width $d_i \in \mathbb{N}$, $i \in I$.

Order:

	Width	Rolls
i	w_i	d_i
1	1380	22
2	1520	25
3	1560	12
4	1710	14
5	1820	18
6	1880	18
7	1930	20
8	2000	10
9	2050	12
10	2100	14
11	2140	16
12	2150	18
13	2200	20

Solution:

	1380	1520	1560	1710	1820	1880	1930	2000	2050	2100	2140	2150	2200
1	1820	1820	1820										
2	1380	2150	1930										
3	1380	2150	2050										
4	1380	2100	2100										
5	2200	1820	1560										
6	2200	1520	1880										
7	1520	1930	2150										
8	1520	1930	2140										
9	1710	2000	1880										
10	1710	1710	2150										

Width of raw paper roll: W

$$W = 5600$$

Customer order:

- ▶ Set I of widths.
- ▶ Widths $w_i \in \mathbb{N}$, $i \in I$.
- ▶ Required number of rolls for each width $d_i \in \mathbb{N}$, $i \in I$.

Order:

	Width	Rolls
i	w_i	d_i
1	1380	22
2	1520	25
3	1560	12
4	1710	14
5	1820	18
6	1880	18
7	1930	20
8	2000	10
9	2050	12
10	2100	14
11	2140	16
12	2150	18
13	2200	20

Set of valid patterns P :

- ▶ $a_{p,i} \in \mathbb{N}$, $p \in P$, $i \in I$: Number of width i in pattern p
- ▶ e.g., $a_{p,\cdot} = (1, 0, 0, 0, 0, 0, 0, 0, 0, 2, 0, 0, 0) \Leftrightarrow 1 \times 1380, 2 \times 2100$
- ▶ pattern only valid if $\sum_{i \in I} a_{p,i} w_i \leq W$.

Solution:

	1380	1520	1560	1710	1820	1880	1930	2000	2050	2100	2140	2150	2200
1380	1820	1820	1820										
1520	1380	2150	1930										
1560	1380	2150	2050										
1710	1380	2100	2100										
1820	2200	1820	1560										
1880	2200	1520	1880										
1930	1520	1930	2150										
2000	1520	1930	2140										
2050	1710	2000	1880										
2100	1710	1710	2150										

Width of raw paper roll: W

$$W = 5600$$

Customer order:

Order:

- ▶ Set I of widths.
- ▶ Widths $w_i \in \mathbb{N}$, $i \in I$.
- ▶ Required number of rolls for each width $d_i \in \mathbb{N}$, $i \in I$.

	Width	Rolls
I	w_i	d_i
1	1380	22
2	1520	25
3	1560	12
4	1710	14
5	1820	18
6	1880	18
7	1930	20
8	2000	10
9	2050	12
10	2100	14
11	2140	16
12	2150	18
13	2200	20

Set of valid patterns P :

- ▶ $a_{p,i} \in \mathbb{N}$, $p \in P$, $i \in I$: Number of width i in pattern p
- ▶ e.g., $a_{p,\cdot} = (1, 0, 0, 0, 0, 0, 0, 0, 0, 2, 0, 0, 0) \Leftrightarrow 1 \times 1380, 2 \times 2100$
- ▶ pattern only valid if $\sum_{i \in I} a_{p,i} w_i \leq W$.

Model:

- ▶ Variable $x_p \in \mathbb{N}$: how often to use pattern p
- ▶ Satisfy demand: $\sum_{p \in P} a_{p,i} x_p \geq d_i$ for all $i \in I$
- ▶ Minimize number of patterns: $\min \sum_{p \in P} x_p$

Solution:

	1380	1520	1560	1710	1820	1880	1930	2000	2050	2100	2140	2150	2200
1380	1820	1820	1820										
1520	1380	2150	1930										
1560	1380	2150	2050										
1710	1380	2100	2100										
1820	2200	1820	1560										
1880	2200	1520	1880										
1930	1520	1930	2150										
2000	1520	1930	2140										
2050	1710	2000	1880										
2100	1710	1710	2150										

Width of raw paper roll: W

$W = 5600$

Customer order:

Order:

- ▶ Set I of widths.
- ▶ Widths $w_i \in \mathbb{N}$, $i \in I$.
- ▶ Required number of rolls for each width $d_i \in \mathbb{N}$, $i \in I$.

	Width	Rolls
i	w_i	d_i
1	1380	22
2	1520	25
3	1560	12
4	1710	14
5	1820	18
6	1880	18
7	1930	20
8	2000	10
9	2050	12
10	2100	14
11	2140	16
12	2150	18
13	2200	20

Set of valid patterns P :

- ▶ $a_{p,i} \in \mathbb{N}$, $p \in P$, $i \in I$: Number of width i in pattern p
- ▶ e.g., $a_{p,\cdot} = (1, 0, 0, 0, 0, 0, 0, 0, 0, 2, 0, 0, 0) \Leftrightarrow 1 \times 1380, 2 \times 2100$
- ▶ pattern only valid if $\sum_{i \in I} a_{p,i} w_i \leq W$.

Model:

- ▶ Variable $x_p \in \mathbb{N}$: how often to use pattern p
- ▶ Satisfy demand: $\sum_{p \in P} a_{p,i} x_p \geq d_i$ for all $i \in I$
- ▶ Minimize number of patterns: $\min \sum_{p \in P} x_p$

Solution:

	1380	1520	1560	1710	1820	1880	1930	2000	2050	2100	2140	2150	2200
1380	1												
1520		1											
1560			1										
1710				1									
1820					1								
1880						1							
1930							1						
2000								1					
2050									1				
2100										1			
2140											1		
2150												1	
2200													1

Observation: Only a few of all valid patterns will be used in a solution \Rightarrow **generate set of active patterns** dynamically

Let $\mathcal{P} \subseteq P$ be the **set of generated pattern**.

Restricted Master Problem:

$$\begin{aligned} \min \quad & \sum_{p \in \mathcal{P}} x_p \\ \text{s.t.} \quad & \sum_{p \in \mathcal{P}} a_{p,i} x_p \geq d_i & \forall i \in I & \quad (\text{RMP}) \\ & x_p \in \mathbb{R}_+ & \forall p \in \mathcal{P} \end{aligned}$$

Let μ be the dual variable associated with the inequality constraint in (RMP).

Pricing Problem:

$$\begin{aligned} \min \quad & 1 - \sum_{i \in I} \mu_i y_i \\ \text{s.t.} \quad & \sum_{i \in I} w_i y_i \leq W & \quad (\text{PP}) \\ & y_i \in \mathbb{N}_+ & \forall i \in I \end{aligned}$$

If (PP) has a solution y^* with value < 0 , an improving pattern has been found \Rightarrow add to \mathcal{P} with $a_{p,i} = y_i^*$.

$$\min \left\{ \sum_{p \in \mathcal{P}} x_p : \sum_{p \in \mathcal{P}} a_{p,i} x_p \geq d_i \forall i \in I, \quad x_p \in \mathbb{R}_+ \right\}$$

```

Set      i      "widths"      / w1*w4 /;
Parameter r      "raw width" / 100 /
          w(i)    "width"      / w1 45, w2 36, w3 31, w4 14 /
          d(i)    "demand"     / w1 97, w2 610, w3 395, w4 211 /;

Set      p      "possible patterns" / 1*1000 /
          pp(p)   "dynamic subset of p";
Parameter
          api(p,i) "number of width i in pattern p";

Variable xp(p)   "how often to use pattern p"
          z       "objective variable"

Integer Variable xp;
xp.up(p) = sum(i, d(i));

Equations numpat      "number of patterns used"
          demand(i)   "meet demand";
numpat..      z =e= sum(pp, xp(pp));
demand(i)..   sum(pp, api(pp,i)*xp(pp)) =g= d(i);

Model master / numpat, demand /;
Solve master min z using RMIP;

```



$$\min \left\{ 1 - \sum_{i \in I} \mu_i y_i : \sum_{i \in I} w_i y_i \leq W, \quad y_i \in \mathbb{N}_+ \right\} \quad (\text{PP})$$

```

Set      i      "widths"      / w1*w4 /;
Parameter r      "raw width" / 100 /
          w(i)    "width"      / w1 45, w2 36, w3 31, w4 14 /;

Variable z      "objective variable"
          y(i)   "new pattern";
Integer variable y;
y.up(i) = ceil(r/w(i));

Equations defobj  "pricing problem objective"
          knapsack "knapsack constraint";

defobj..      z =e= 1 - sum(i, demdual(i)*y(i));
knapsack..    sum(i, w(i)*y(i)) =l= r;

Model pricing / defobj, knapsack /;
Solve pricing min z using MIP;

```


Task: Get the Python file `cutstock.py` and add the Column-Generation loop:

1. Solve (RMP) as LP.
2. Update μ in (PP) and solve (PP).
3. If optimal value of (PP) is > 0 , then construct new pattern for (RMP) from solution of (PP) and go to 1.

Final master objective value should be **452.25**.

Task: Get the Python file `cutstock.py` and add the Column-Generation loop:

1. Solve (RMP) as LP.
2. Update μ in (PP) and solve (PP).
3. If optimal value of (PP) is > 0 , then construct new pattern for (RMP) from solution of (PP) and go to 1.

Final master objective value should be **452.25**.

Solution File: `cutstock_sol.py`