

# The MIP Landscape

# Example 1: LP still can be HARD

SGM: Schedule Generation Model

157323 rows, 182812 columns, 6348437 nzs

- ❑ LP relaxation at root node:
  - 18 hours
- ❑ Branch-and-bound
  - 1710 nodes, first feasible
  - 3.7% gap
  - Time: **92 days!!**
- ❑ MIP does not appear to be difficult: *LP is a roadblock*

# Example 2: MIP really is HARD

A customer model: 44 cons, 51 vars, 167 nzs, maximization  
51 general integer variables (*and no bounds*)

Branch-and-bound: Initial integer solution    -2186.0  
                          Initial upper bound       -1379.4

...after 1.4 days, 32,000,000 B&B nodes, 5.5 Gig tree

Integer solution and bound: **UNCHANGED**

**What's wrong?** Bad modeling. Free GIs chase each other off to infinity.

## Example 2 (cont.): Here's what's wrong

```
Maximize
  x + y + z
Subject To
  2 x + 2 y ≤ 1
  z = 0
  x free y free
  x,y integer
```

Note: This problem can be solved in several ways

- Removing  $z=0$ , objective is integral [*Presolve*]
- Euclidean reduction on the constraint [*Presolve*]

However: Branch-and-bound cannot solve!

# Example 3: A typical situation today

## – Supply-chain scheduling

- Model description:
  - Weekly model, daily buckets: Objective to minimize end-of-day inventory.
  - Production (single facility), inventory, shipping (trucks), wholesalers (demand known)
- Initial modeling phase
  - Simplified prototype + complicating constraints (production run grouping req't, min truck constraints)
  - **RESULT: Couldn't get good feasible solutions.**
- Decomposition approach
  - Talk to current scheduling team: They first decide on “producibles” schedule. Simulate using heuristics.
  - **Fixed model: Fix variables and run MIP**

# Supply-chain scheduling (continued): Solving the fixed model

## CPLEX 5.0:

```
Integer optimal solution (0.0001/0): Objective = 1.5091900536e+05  
Current MIP best bound = 1.5090391809e+05 (gap = 15.0873)  
Solution time = 3465.73 sec. Iterations = 7885711 Nodes = 489870 (2268)
```

## Gurobi 2.0:

Cutting planes:

Gomory: 16

Implied bound: 33

MIR: 1

Flow cover: 70

Explored 9 nodes (4517 simplex iterations) in 0.49 seconds

Optimal solution found (tolerance 1.00e-04)

Best objective 1.5091900536e+05, best bound 1.5090900608e+05, gap 0.0066%

**Original model:** Now solves to optimality in 1  
hour (20% improvement in solution quality)

# Example 4: Unit-Commitment Model

Electrical Power Industry, ERPI GS-6401, June 1989:  
Mixed-integer programming (MIP) is a powerful modeling tool, “They are, however, theoretically complicated and computationally cumbersome”

*In Other Words:* MIP is an interesting “toy”, but it just isn’t going to work in practice.

# California 7-Day Model

**UNITCAL\_7**: 48939 constraints, 25755 variables (2856 binary)

**Reported Results 1999** – machine unknown

2 Day model: 8 hours, no progress

7 Day model: 1 hour to solve initial LP

**Desktop PC** -- ran full 7-day model

**CPLEX 6.5 (1999)**: 22 minutes, optimal



# California 7-Day Model

Gurobi Optimizer version 2.0.0

Read MPS format model from file unitcal\_7.mps.bz2  
Optimize a model with 48939 Rows, 25755 Columns and 127595 NonZeros  
Presolved: 38804 Rows, 19960 Columns, 105627 Nonzeros

Root relaxation: objective 1.945018e+07, 18340 iterations, 0.60 seconds

Nodes		Current Node			Objective Bounds			Work		
Expl	Unexpl	Obj	Depth	IntInf	Incumbent	BestBd	Gap	It/Node	Time	
0	0	1.9450e+07	0	772	-	1.9450e+07	-	-	2s	
0	0	1.9582e+07	0	673	-	1.9582e+07	-	-	8s	
0	0	1.9590e+07	0	543	-	1.9590e+07	-	-	13s	
0	0	1.9592e+07	0	530	-	1.9592e+07	-	-	15s	
0	0	1.9592e+07	0	521	-	1.9592e+07	-	-	16s	
0	0	1.9592e+07	0	531	-	1.9592e+07	-	-	16s	
H	0	0			2.0669e+07	1.9592e+07	5.21%	-	18s	
	0	2	1.9592e+07	0	531	2.0669e+07	1.9592e+07	5.21%	-	19s
	3	3	1.9595e+07	2	581	2.0669e+07	1.9593e+07	5.20%	567	20s
	32	35	1.9665e+07	8	394	2.0669e+07	1.9605e+07	5.15%	1043	30s
H	94	92			1.9849e+07	1.9605e+07	1.23%	565	35s	
H	122	51			1.9646e+07	1.9605e+07	0.21%	477	39s	
	385	56	1.9638e+07	31	33	1.9646e+07	1.9605e+07	0.21%	200	42s
H	747	167			1.9636e+07	1.9608e+07	0.14%	146	48s	
H	818	199			1.9636e+07	1.9610e+07	0.13%	142	52s	
	936	248	1.9630e+07	21	113	1.9636e+07	1.9611e+07	0.13%	137	57s
H	1053	304			1.9636e+07	1.9611e+07	0.12%	130	60s	
	1182	379	1.9622e+07	16	531	1.9636e+07	1.9612e+07	0.12%	124	70s
	1193	386	1.9632e+07	26	520	1.9636e+07	1.9612e+07	0.12%	122	90s
	1199	390	1.9623e+07	18	520	1.9636e+07	1.9612e+07	0.12%	122	95s
	1208	398	1.9612e+07	15	640	1.9636e+07	1.9612e+07	0.12%	157	199s
	1209	399	1.9612e+07	15	614	1.9636e+07	1.9612e+07	0.12%	159	205s
	1285	370	1.9635e+07	27	525	1.9636e+07	1.9613e+07	0.11%	169	227s
	2266	300	1.9634e+07	33	329	1.9636e+07	1.9633e+07	0.01%	113	235s

Cutting planes:

Gomory: 32  
Cover: 22  
Implied bound: 599  
Clique: 60  
Flow cover: 391

Explored 2379 nodes (298817 simplex iterations) in 235.33 seconds

Optimal solution found (tolerance 1.00e-04)

Best objective 1.9635558244e+07, best bound 1.9633702475e+07, gap 0.0095%

# Computational History: 1950 – 1998

- **1954 Dantzig, Fulkerson, S. Johnson:** 42 city TSP
  - Solved to optimality using LP and cutting planes
- **1957 Gomory**
  - Cutting plane algorithms
- **1960 Land, Doig, 1965 Dakin**
  - B&B
- **1971 MPSX/370**
- **1972 UMPIRE**
  - LP-based B&B
  - MIP became commercially viable
- **1972 – 1998** Good B&B remained the state-of-the-art in commercial codes, in spite of ....
  - Edmonds, polyhedral combinatorics
  - 1973 Padberg, cutting planes
  - 1973 Chvátal, revisited Gomory
  - 1974 Balas, disjunctive programming
  - 1983 Crowder, Johnson, Padberg: PIPX, pure 0/1 MIP
  - 1987 Van Roy and Wolsey: MPSARX, mixed 0/1 MIP
  - TSP, Grötschel, Padberg, ...

# 1998 ... A New Generation of MIP Codes

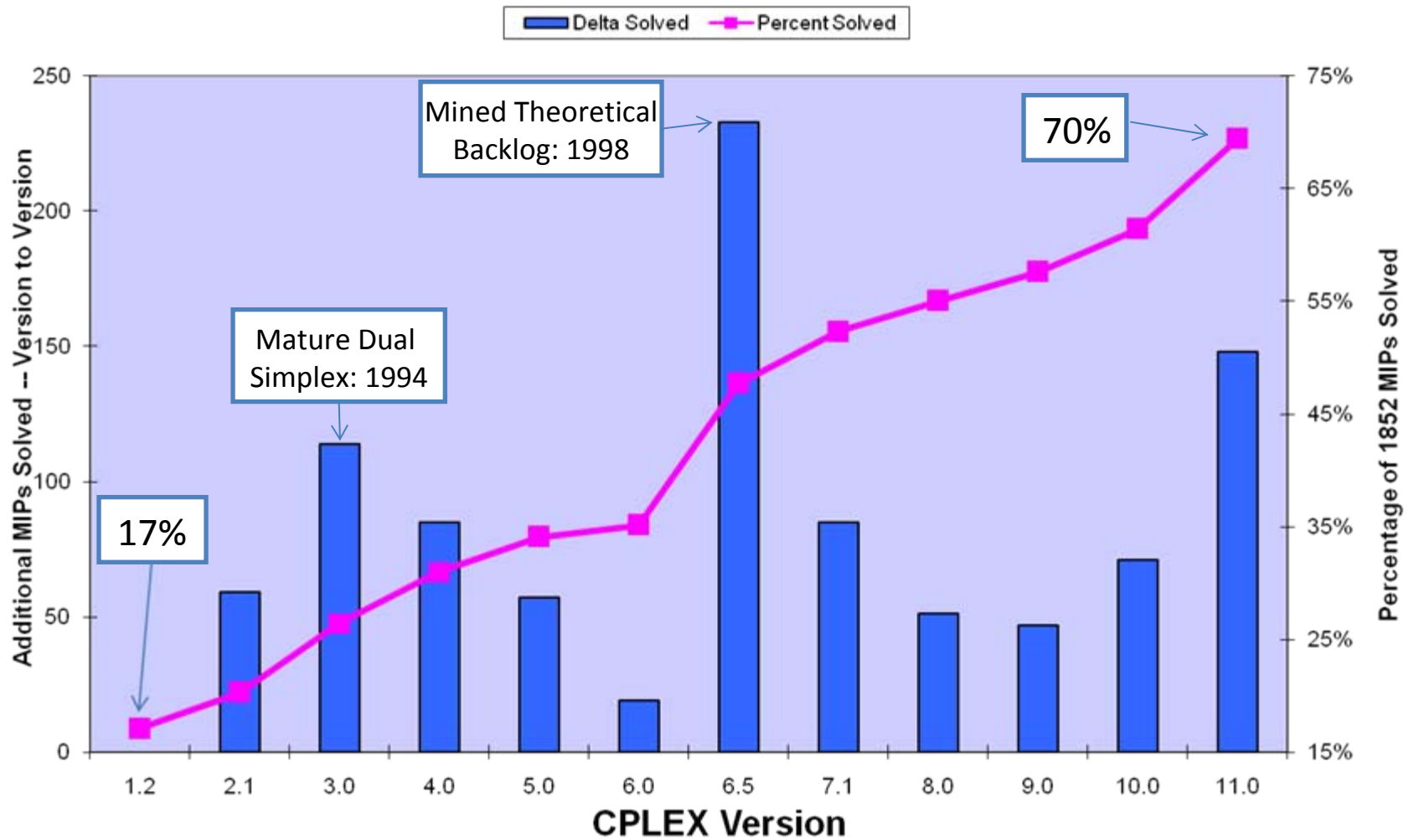
- Linear programming
  - Stable, robust dual simplex
- Variable/node selection
  - Influenced by traveling salesman problem
- Primal heuristics
  - 12 different tried at root
  - Retried based upon success
- Node presolve
  - Fast, incremental bound strengthening (very similar to Constraint Programming)
- Presolve – numerous small ideas
  - Probing in constraints:  
$$\sum x_j \leq (\sum u_j) y, \quad y = 0/1$$
$$\rightarrow x_j \leq u_j y \text{ (for all } j)$$
- Cutting planes
  - Gomory, mixed-integer rounding (MIR), knapsack covers, flow covers, cliques, GUB covers, implied bounds, zero-half cuts, path cuts

# Some Test Results

- **Test set: 1852 real-world MIPs**
  - Full library
    - 2791 MIPs
  - Removed:
    - 559 “Easy” MIPs
    - 348 “Duplicates”
    - 22 “Hard” LPs (0.8%)
- **Parameter settings**
  - Pure defaults
  - 30000 second time limit
- **Versions Run**
  - CPLEX 1.2 (1991) -- CPLEX 11.0 (2007)

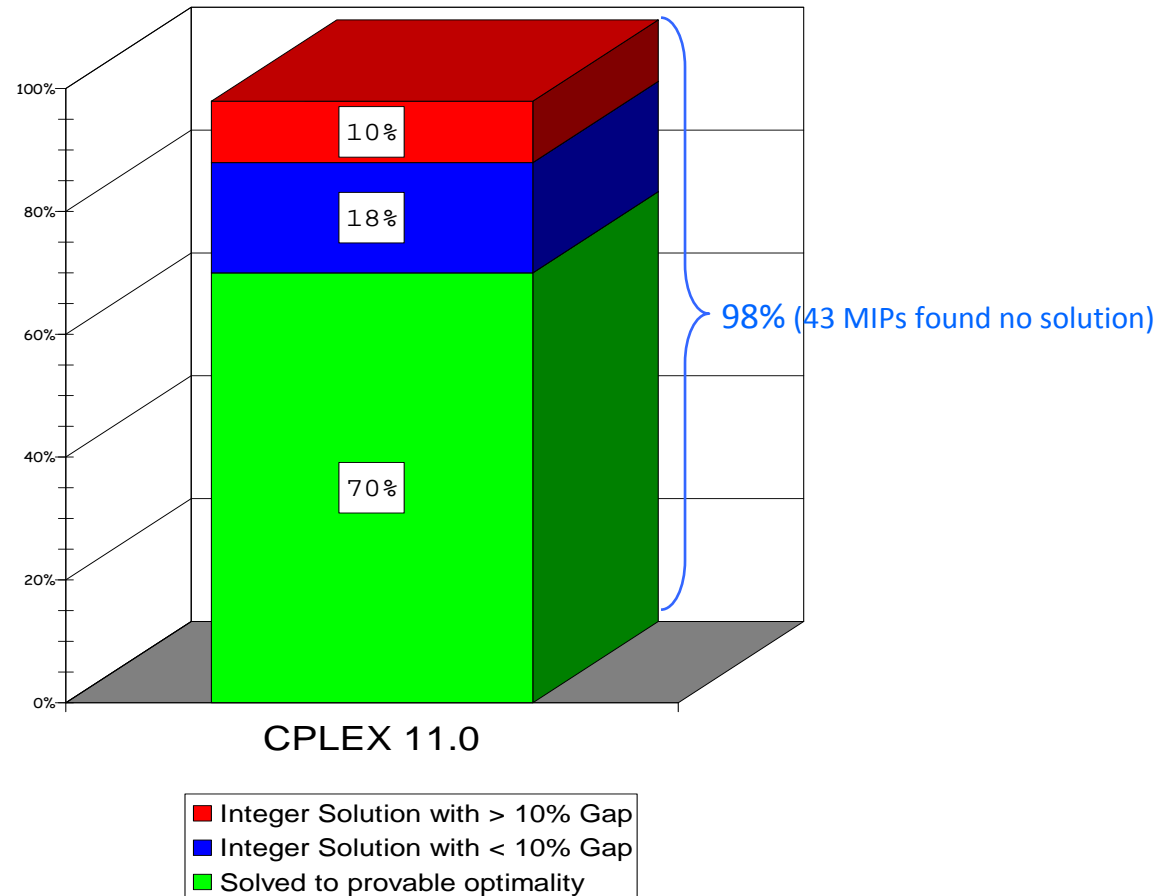
# MIP Solvability

# Solvability of MIPS 1991-Present



# Solvability of MIPs – CPLEX 11.0

1852 MIPs, 30000 second time limit



# MIP Speedups



# Speedups 1991-Present

